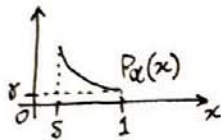


ESTIMATION OF BRADLEY-TERRY MODEL PARAMETERS: (with Ali Jadbabaie & Devavrat Shah)① MOTIVATION:

- Measure the level of skill in sports
 - [Getty et al 2018]: Luck vs Skill in fantasy sports ← motivated by policy-making on gambling
 - [Misra-Shah-Ranganathan 2020]: Hypothesis testing for pure skill
- Statistical Formulation:
 - Fix constants $\delta \in (0, 1)$ and $\sigma > 0$.
 - Let $\mathcal{P}_{\delta, \sigma}(\sigma) =$ set of all probability density functions (PDFs) on $[s, 1]$ that are $\geq \delta$ and σ -Lipschitz continuous, i.e. $P_\alpha \in \mathcal{P}_{\delta, \sigma}(\sigma) \Leftrightarrow P_\alpha$ is a PDF and $\forall x, y \in [s, 1], |P_\alpha(x) - P_\alpha(y)| \leq \sigma|x - y|, \forall x \in [s, 1], P_\alpha(x) \geq \delta$.
 - Each sport has an unknown PDF $P_\alpha \in \mathcal{P}_{\delta, \sigma}(\sigma)$ of merit values.

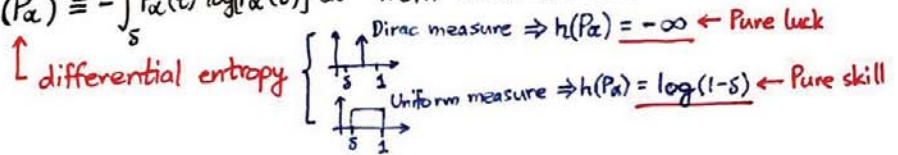


- Suppose there is a tournament with $n \geq 2$ players: $\{1, \dots, n\}$. Each player i has merit value $\alpha_i \sim P_\alpha$ so that $\alpha_1, \dots, \alpha_n \stackrel{iid}{\sim} P_\alpha$.
- There are $\binom{n}{2}$ independent two-player games in a tournament, with likelihoods:

$$\mathbb{P}(j \text{ beats } i | \alpha_1, \dots, \alpha_n) = \frac{\alpha_j}{\alpha_i + \alpha_j} \text{ for all } i \neq j.$$

← scale invariance

This is the Bradley-Terry Model!
- We see observations: $\{Z(i, j) \triangleq \mathbb{1}\{j \text{ beats } i\} : 1 \leq i < j \leq n\}$. (For $i > j$, $Z(i, j) = 1 - Z(j, i)$, and $Z(i, i) = 0$.)
- GOAL: Estimate P_α or $h(P_\alpha) \triangleq -\int_s^1 P_\alpha(t) \log[P_\alpha(t)] dt$ from observations.

• Approach:

- Focus of Talk**
- 1) Estimate Bradley-Terry Model parameters $\alpha_1, \dots, \alpha_n$ based on observations. Suppose estimates are $\hat{\alpha}_1, \dots, \hat{\alpha}_n$.
 - 2) Estimate P_α or $h(P_\alpha)$ using $\hat{\alpha}_1, \dots, \hat{\alpha}_n$.
(Robust kernel density estimation ← use $\hat{\alpha}_1, \dots, \hat{\alpha}_n$ instead of $\alpha_1, \dots, \alpha_n$.)
- $\hookrightarrow \hat{P}_\alpha(x) = \frac{1}{n} \sum_{i=1}^n \frac{1}{h} K\left(\frac{x - \hat{\alpha}_i}{h}\right)$
- ↑ kernel bandwidth
Parzen-Rosenblatt estimator

② BRADLEY-TERRY MODEL: MINIMAX ESTIMATION

- [Bradley-Terry 1952] (originally proposed by [Zermelo 1929]): Ranking based on paired comparisons.
 - n items $\{1, \dots, n\}$ with underlying merits $\alpha_1, \dots, \alpha_n \geq 0$
 - Easy to compare any two, but hard to rank all.
 - Use model $\mathbb{P}(i > j) = \frac{\alpha_i}{\alpha_i + \alpha_j}$ of pairwise comparisons to find "true" merits $\alpha_1, \dots, \alpha_n$.

[continued.]

• Plackett-Luce Model: [Luce 1959], [Plackett 1975] ← social choice theory/econometrics

- Luce's choice axiom: Probability of selecting one item over another in a set of items is not affected by the presence or absence of other items in the set.
 ↑ independence of irrelevant alternatives ↑ axiom for prob. model over selection

- Equivalent model: n items $\{1, \dots, n\}$ with merits $\alpha_1, \dots, \alpha_n \geq 0$

$$P(\text{select } i) = \frac{\alpha_i}{\sum_{j=1}^n \alpha_j}$$

- Distribution over permutations/rankings:

$$\forall \sigma, \quad P_L(\sigma) = \frac{\alpha_{\sigma(1)}}{\sum_{k \in \{1, \dots, n\}} \alpha_k} \cdot \frac{\alpha_{\sigma(2)}}{\sum_{k \in \{1, \dots, n\} \setminus \{\sigma(1)\}} \alpha_k} \cdot \frac{\alpha_{\sigma(3)}}{\sum_{k \in \{1, \dots, n\} \setminus \{\sigma(1), \sigma(2)\}} \alpha_k} \cdots \frac{\alpha_{\sigma(n)}}{\alpha_{\sigma(n)}} \quad \left. \vphantom{\frac{\alpha_{\sigma(1)}}{\sum_{k \in \{1, \dots, n\}} \alpha_k}} \right\} \text{Plackett-Luce model}$$

↑ permutation of $\{1, \dots, n\}$

- Pairwise selection → Bradley-Terry model

• Thurstonian Model: [Thurstone 1927] ← psychometrics

- Law of Comparative Judgment: "Discriminal" process to rank n items $\{1, \dots, n\}$ is modeled by first associating merits $\alpha_1, \dots, \alpha_n \geq 0$ to the items, and then ranking them by ranking the n random variables $\alpha_1 + X_1, \dots, \alpha_n + X_n$ for i.i.d. X_1, \dots, X_n .
 noise in discriminial process

- Distribution over permutations/rankings:

$$\forall \sigma, \quad P_T(\sigma) = P(\alpha_{\sigma(1)} + X_{\sigma(1)} > \alpha_{\sigma(2)} + X_{\sigma(2)} > \dots > \alpha_{\sigma(n)} + X_{\sigma(n)})$$

↑ permutation of $\{1, \dots, n\}$

- Equivalent to Plackett-Luce model if and only if $X_1, \dots, X_n \stackrel{\text{iid}}{\sim} \text{Gumbel}(\mu, \beta)$. [Yellott 1977]
 generalized extreme value (Type I) dist.
 ↑ real location ↑ positive scale

(Note: $F_X(x) \triangleq e^{-e^{-(x-\mu)/\beta}}$ is CDF of Gumbel dist.)

• Minimax Formulation:

Define $\pi(i) \triangleq \frac{\alpha_i}{\sum_{j=1}^n \alpha_j}$, $\forall i \in \{1, \dots, n\}$, and let $\pi \triangleq (\pi(1), \dots, \pi(n))$.
 ↑ canonically scaled merit parameters

Find upper and lower bounds on:

$$\inf_{\hat{\pi}} \sup_{P \in \mathcal{P}(\sigma)} E[\|\hat{\pi} - \pi\|_1]$$

↑ ℓ^1 -norm ↑ wrt $P_{\alpha}^{\text{PL}}, P_{\alpha_1, \dots, \alpha_n}$

where infimum is over all estimators $\hat{\pi}$ of π based on $\{Z(i, j) : i < j\}$.

↑ could be randomized, must be a prob. mass function

[continued.]

③ MINIMAX UPPER BOUND: ← construct estimator

• Rank Centrality: [Negahban-Oh-Shah 2017]

- Define the row stochastic matrix $S \in \mathbb{R}^{n \times n}$:

$$\forall i \neq j, S(i, j) \triangleq \frac{1}{n-1} \cdot \frac{\alpha_j}{\alpha_i + \alpha_j} > 0, \quad \text{since } \alpha_i's \geq \delta$$

\uparrow (i, j) th entry of S

$$\forall i, S(i, i) \triangleq 1 - \frac{1}{n-1} \sum_{k \neq i} \frac{\alpha_k}{\alpha_i + \alpha_k} = \frac{1}{n-1} \sum_{k \neq i} \frac{\alpha_i}{\alpha_i + \alpha_k} > 0, \quad \text{since } \alpha_i's \geq \delta$$

(Clearly, $\sum_{j=1}^n S(i, j) = 1, \forall i$.)

- S defines a Markov chain on the state-space of players $\{1, \dots, n\}$.

- Detailed Balance Conditions:

$$\forall i \neq j, \pi(i) S(i, j) = \frac{\alpha_i}{\sum_{k=1}^n \alpha_k} \cdot \frac{1}{n-1} \cdot \frac{\alpha_j}{\alpha_i + \alpha_j} = \frac{\alpha_j}{\sum_{k=1}^n \alpha_k} \cdot \frac{1}{n-1} \cdot \frac{\alpha_i}{\alpha_i + \alpha_j} = \pi(j) S(j, i).$$

S self-adjoint operator
 S satisfies Kolmogorov criterion

Hence, S defines a reversible Markov chain with invariant distribution π : $\pi = \pi S$.

- π is unique as $S > 0$ entry-wise $\Rightarrow S$ ergodic, i.e. irreducible & aperiodic.

- Construct estimator $\tilde{S} \in \mathbb{R}^{n \times n}$ of S based on $Z(i, j)$'s:

$$\forall i \neq j, \tilde{S}(i, j) \triangleq \frac{1}{n-1} Z(i, j) \geq 0,$$

$$\forall i, \tilde{S}(i, i) \triangleq 1 - \frac{1}{n-1} \sum_{k \neq i} Z(i, k) = \frac{1}{n-1} \sum_{k \neq i} Z(k, i) \geq 0.$$

(Clearly, $\sum_{j=1}^n \tilde{S}(i, j) = 1, \forall i$.)

\tilde{S} is row stochastic

- \tilde{S} defines another Markov chain. ← not necessarily reversible or ergodic
Let $\tilde{\pi} = \tilde{\pi} \tilde{S}$ be any invariant distribution of \tilde{S} .

- $\tilde{\pi}$ is an estimator of π .

Thm: [Chen et al 2019]

$$a) \frac{\|\tilde{\pi} - \pi\|_{\infty}}{\|\pi\|_{\infty}} = O\left(\frac{1}{\delta} \sqrt{\frac{\log(n)}{n}}\right) \text{ with probability } \geq 1 - O(n^{-5}).$$

$$b) \frac{\|\tilde{\pi} - \pi\|_2}{\|\pi\|_2} = O\left(\frac{1}{\sqrt{n}}\right) \text{ with probability } \geq 1 - O(n^{-5}).$$

\uparrow if $\delta = \Theta(1)$

- Can we bound $\|\tilde{\pi} - \pi\|_1$ with high probability?

Yes!

[continued.]

Theorem: $\inf_{\hat{\pi}} \sup_{P_{\alpha} \in \mathcal{P}_{[8,1]}(\sigma)} \mathbb{E}[\|\hat{\pi} - \pi\|_1] \leq \sup_{P_{\alpha} \in \mathcal{P}_{[8,1]}(\sigma)} \mathbb{E}[\|\hat{\pi} - \pi\|_1] = O\left(\frac{1}{\sqrt{n}}\right)$
 (Upper Bound)
 for all sufficiently large n .

- Proof:

Using [Chen et al 2019] part (b), with probability $\geq 1 - O(n^{-5})$,

$$\|\hat{\pi} - \pi\|_1 \leq \sqrt{n} \|\hat{\pi} - \pi\|_2 \leq \overset{\text{some constant}}{C} \|\pi\|_2$$

for all sufficiently large n , using equivalence of norms.

Since $\|\pi\|_2 \leq \frac{1}{\sqrt{n}}$ (because $\alpha_i \in [8, 1]$), we get:
 \uparrow constant

$$\|\hat{\pi} - \pi\|_1 = O\left(\frac{1}{\sqrt{n}}\right) \text{ with probability } \geq 1 - O(n^{-5}).$$

The law of total expectation and the bound $\|\hat{\pi} - \pi\|_1 \leq \|\hat{\pi}\|_1 + \|\pi\|_1 = 2$ (triangle inequality) yield the desired result. ▀

④ MINIMAX LOWER BOUND:

• Bayes Risk:

$$\inf_{\hat{\pi}} \sup_{P_{\alpha} \in \mathcal{P}_{[8,1]}(\sigma)} \mathbb{E}[\|\hat{\pi} - \pi\|_1] \geq \overset{\text{Bayes risk}}{\inf_{\hat{\pi}} \mathbb{E}[\|\hat{\pi} - \pi\|_1]}$$

\uparrow Choose $P_{\alpha} = \text{Uniform}([8, 1])$.

How do we lower bound Bayes risk?

• Generalized Fano's Method:

$$\text{Lemma: [Xu-Raginsky 2017]} \quad \inf_{\hat{\pi}} \mathbb{E}[\|\hat{\pi} - \pi\|_1] \geq \sup_{t > 0} t \left(1 - \frac{I(\alpha^n; Z) + \log(Z)}{\log(1/\mathcal{L}(t))} \right)$$

$\alpha^n \triangleq (\alpha_1, \dots, \alpha_n)$ $Z \triangleq \{Z(i, j) : i < j\}$

where $I(\alpha^n; Z) \triangleq D(P_{\alpha^n, Z} \| P_{\alpha^n} P_Z)$ and $\mathcal{L}(t) \triangleq \sup_{\nu} \mathbb{P}(\|\pi - \nu\|_1 \leq t)$ for $t > 0$.

\uparrow mutual information

\uparrow KL divergence

\uparrow small ball probability (measure of concentration)

- Intuition: $\mathcal{L}(t) \uparrow \Rightarrow$ high conc. of $\pi \Rightarrow$ Bayes risk \downarrow (as we can estimate π easily).

$I(\alpha^n; Z) \uparrow \Rightarrow Z$ has lots of info. about $\pi \Rightarrow$ Bayes risk \downarrow .

- Proof: Fix any $\hat{\pi}$ and any $t > 0$.

Consider Markov chain $\pi \rightarrow \alpha^n \rightarrow Z \rightarrow \hat{\pi}$.

$$I(\alpha^n; Z) \geq I(\pi; \hat{\pi}) \quad [\text{DPI}]$$

$$= D(P_{\pi, \hat{\pi}} \| P_{\pi} \cdot P_{\hat{\pi}}) = D(P_{\pi, \hat{\pi}} \| Q_{\pi, \hat{\pi}}), \text{ where } Q_{\pi, \hat{\pi}} \triangleq P_{\pi} \cdot P_{\hat{\pi}}$$

$$\geq D(P_{\pi, \hat{\pi}}(\|\pi - \hat{\pi}\|_1 \leq t) \| Q_{\pi, \hat{\pi}}(\|\pi - \hat{\pi}\|_1 \leq t)) \quad [\text{DPI: } (\pi, \hat{\pi}) \mapsto \mathbb{1}\{\|\hat{\pi} - \pi\|_1 \leq t\}]$$

↑ binary KL divergence

$$= D(P_{\pi, \hat{\pi}}(\|\pi - \hat{\pi}\|_1 \leq t) \| \mathbb{E}_{\hat{\pi}}[P_{\pi}(\|\pi - \hat{\pi}\|_1 \leq t)])$$

$$\geq P_{\pi, \hat{\pi}}(\|\pi - \hat{\pi}\|_1 \leq t) \log \left(\frac{1}{\mathbb{E}_{\hat{\pi}}[P_{\pi}(\|\pi - \hat{\pi}\|_1 \leq t)]} \right) - \log(2)$$

$$\geq P(\|\pi - \hat{\pi}\|_1 \leq t) \log \left(\frac{1}{\mathcal{L}(t)} \right) - \log(2).$$

$$\Rightarrow P(\|\hat{\pi} - \pi\|_1 > t) \geq 1 - \frac{I(\alpha^n; Z) + \log(2)}{\log \left(\frac{1}{\mathcal{L}(t)} \right)}.$$

Lemma: $\forall p, q \in [0, 1]$,

$$\begin{aligned} D(p \| q) &= p \log \left(\frac{p}{q} \right) + (1-p) \log \left(\frac{1-p}{1-q} \right) \\ &= p \log \left(\frac{p}{q} \right) + (1-p) \log \left(\frac{1-p}{1-q} \right) \\ &\quad \leftarrow \text{binary entropy} \\ &\geq p \log \left(\frac{1}{q} \right) - \log(2). \end{aligned}$$

By Markov's inequality:

$$\begin{aligned} \mathbb{E}[\|\hat{\pi} - \pi\|_1] &\geq t P(\|\hat{\pi} - \pi\|_1 > t) \\ &\geq t \left(1 - \frac{I(\alpha^n; Z) + \log(2)}{\log \left(\frac{1}{\mathcal{L}(t)} \right)} \right). \end{aligned}$$

Take $\inf_{\hat{\pi}}$ and $\sup_{t>0}$. This completes the proof. ■

• Bound $I(\alpha^n; Z)$: (Covering Number Method)

- For $\beta = (\beta_1, \dots, \beta_n) \in [s, 1]^n$, let $P_{\beta|B} = P_{Z|\alpha^n=B}$.

- Def: For any $\epsilon > 0$, we say that $\{\beta^{(1)}, \dots, \beta^{(M)}\} \subset [s, 1]^n$ is an ϵ -covering of $[s, 1]^n$ with cardinality M if $\forall \beta \in [s, 1]^n$, $\exists i \in \{1, \dots, M\}$, $D(P_{\beta|B} \| P_{\beta^{(i)}|B}) \leq \epsilon$.

$\Rightarrow M^*(\epsilon) \triangleq \min\{M : \exists \epsilon\text{-covering with cardinality } M\}$.

↑ ϵ -covering number

- Lemma: [Yang-Barron 1999] $I(\alpha^n; Z) \leq \inf_{\epsilon > 0} \epsilon + \log(M^*(\epsilon))$.

Proof: Fix any $\epsilon > 0$. Let $\{\beta^{(1)}, \dots, \beta^{(M^*(\epsilon))}\} \subset [s, 1]^n$ be an ϵ -covering. Let $i(\beta) \triangleq \arg \min_{i \in \{1, \dots, M^*(\epsilon)\}} D(P_{\beta|B} \| P_{\beta^{(i)}|B})$, $\forall \beta \in [s, 1]^n$. ■

$$\begin{aligned} I(\alpha^n; Z) &= D(P_{\alpha^n, Z} \| P_{\alpha^n} \cdot P_Z) \\ &= \mathbb{E}_{P_{\alpha^n, Z}} \left[\log \left(\frac{P_{\alpha^n, Z}}{P_{\alpha^n} \cdot P_Z} \right) \right] = \mathbb{E} \left[\log \left(\frac{P_{Z|\alpha^n} \cdot \frac{1}{M^*} \sum_{i=1}^{M^*} P_{Z|\beta^{(i)}}}{P_Z \cdot \frac{1}{M^*} \sum_{i=1}^{M^*} P_{Z|\beta^{(i)}}} \right) \right] \\ &= \mathbb{E}_{\alpha^n} \left[D(P_{Z|\alpha^n} \| \frac{1}{M^*} \sum_{i=1}^{M^*} P_{Z|\beta^{(i)}}) \right] - \underbrace{D(P_Z \| \frac{1}{M^*} \sum_{i=1}^{M^*} P_{Z|\beta^{(i)}})}_{\geq 0 \text{ [Gibbs]}} \\ &\leq \mathbb{E}_{\alpha^n} \left[\mathbb{E}_{P_{Z|\alpha^n}} \left[\log \left(\frac{P_{Z|\alpha^n}}{\frac{1}{M^*} \sum_{i=1}^{M^*} P_{Z|\beta^{(i)}(\alpha^n)}} \right) \right] \right] \\ &= \mathbb{E}_{\alpha^n} \left[\underbrace{D(P_{Z|\alpha^n} \| P_{Z|\beta^{(i(\alpha^n))}})}_{\leq \epsilon \text{ [}\epsilon\text{-covering]}} \right] + \log(M^*(\epsilon)) \\ &\leq \epsilon + \log(M^*(\epsilon)). \end{aligned}$$

Take $\inf_{\epsilon > 0}$. (Note: Can also take \sup_{α^n} .)

↑ bound on channel capacity

$$\text{Lemma: } I(\alpha^n; Z) \leq \frac{1}{2} n \log(n) + \frac{(1-\delta)^2}{8\delta^2} \left(2 + \delta + \frac{1}{\delta}\right) n.$$

$$\text{Proof: Let } Q \triangleq \left\{ s + \frac{(1-\delta)k}{\sqrt{n}} : k \in \{1, \dots, \lfloor \sqrt{n} \rfloor\} \right\} \leftarrow \text{quantize } [\delta, 1]$$

$$\text{Then, } \forall t \in [\delta, 1], \min_{s \in Q} |t - s| \leq \frac{1-\delta}{\sqrt{n}} \leftarrow \text{floor function}$$

$$\text{Claim: } Q^n \text{ is an } \varepsilon\text{-covering with } |Q^n| = |Q|^n \leq n^{n/2} \text{ and } \varepsilon = \frac{(1-\delta)^2}{8\delta^2} \left(2 + \delta + \frac{1}{\delta}\right) n.$$

$$\rightarrow \text{Pf: Fix any } \beta = (\beta_1, \dots, \beta_n) \in [\delta, 1]^n, \text{ and choose } \gamma = (\gamma_1, \dots, \gamma_n) \in Q^n \text{ such that } |\beta_i - \gamma_i| \leq \frac{1-\delta}{\sqrt{n}}, \forall i.$$

$$\begin{aligned} D(P_{Z|\beta} \| P_{Z|\gamma}) &= \sum_{i < j} D(P_{Z(i,j)|\beta} \| P_{Z(i,j)|\gamma}) \\ &= \sum_{i < j} D\left(\frac{\beta_i}{\beta_i + \beta_j} \parallel \frac{\gamma_i}{\gamma_i + \gamma_j}\right) \leftarrow \text{binary KL divergence} \\ &\leq \sum_{i < j} \chi^2\left(\frac{\beta_i}{\beta_i + \beta_j} \parallel \frac{\gamma_i}{\gamma_i + \gamma_j}\right) \leftarrow \text{binary } \chi^2\text{-divergence} \\ &= \sum_{i < j} \left(\frac{\beta_i}{\beta_i + \beta_j} - \frac{\gamma_i}{\gamma_i + \gamma_j}\right)^2 \left(2 + \frac{\gamma_i}{\gamma_j} + \frac{\gamma_j}{\gamma_i}\right) \\ &\leq \underbrace{\left(2 + \delta + \frac{1}{\delta}\right)}_{\substack{[\delta, \frac{1}{\delta}] \ni t \mapsto t + \frac{1}{t} \\ \text{maximized at endpoints}}} \sum_{i < j} \underbrace{\left(\left|\frac{\beta_i}{\beta_i + \beta_j} - \frac{\gamma_i}{\gamma_i + \gamma_j}\right| + \left|\frac{\beta_j}{\gamma_i + \beta_j} - \frac{\gamma_j}{\gamma_i + \gamma_j}\right|\right)^2}_{\Delta\text{-inequality}} \\ &\leq \left(\frac{1}{4\delta}\right)^2 \left(2 + \delta + \frac{1}{\delta}\right) \sum_{i < j} \underbrace{(|\beta_i - \gamma_i| + |\beta_j - \gamma_j|)^2}_{F: [\delta, \infty)^2 \rightarrow \mathbb{R}, F(x, y) \triangleq \frac{x}{x+y}} \\ &\leq \frac{(1-\delta)^2}{8\delta^2} \left(2 + \delta + \frac{1}{\delta}\right) (n-1) \quad \begin{array}{l} \text{For fixed } x, F \text{ is } \frac{1}{4\delta}\text{-Lipschitz in } y. \\ \text{For fixed } y, F \text{ is } \frac{1}{4\delta}\text{-Lipschitz in } x. \end{array} \\ &\leq \varepsilon. \end{aligned}$$

Using Yang-Barron Lemma,

$$\begin{aligned} I(\alpha^n; Z) &\leq \varepsilon + \log(M^*(\varepsilon)) \leftarrow \text{our } \varepsilon \text{ value} \\ &\leq \varepsilon + \log(|Q^n|) \leftarrow \text{our } \varepsilon\text{-covering} \\ &\leq \frac{1}{2} n \log(n) + \frac{(1-\delta)^2}{8\delta^2} \left(2 + \delta + \frac{1}{\delta}\right) n. \end{aligned}$$

This completes the proof. \square

- Remark: This is better than standard information inequalities (tensorization bounds), which give $I(\alpha^n; Z) = O(n^2)$.

[continued.]

• Bound Small Ball Probability: ← no standard approach in the literature

- Lemma: $\forall t > 0, \mathcal{L}(t) \leq \left(\frac{2e}{1-\delta}\right)^n t^{n-1}$.

Proof: For any $t > 0$,

$$\mathcal{L}(t) = \sup_{\nu} \mathbb{P}(\|\pi - \nu\|_1 \leq t)$$

$$\leq \sup_{\tilde{\nu}} \mathbb{P}(\|\tilde{\pi} - \tilde{\nu}\|_1 \leq t) \leftarrow \begin{array}{l} \tilde{\pi} = (\pi(1), \dots, \pi(n-1)) \\ \tilde{\nu} = (\nu_1, \dots, \nu_{n-1}) \text{ where } \nu = (\nu_1, \dots, \nu_n) \end{array}$$

$$= \sup_{\tilde{\nu}} \int_{\mathbb{B}^{n-1}} \underbrace{P_{\tilde{\pi}}(\tau)}_{\text{PDF of } \tilde{\pi} \text{ (Note: } \pi \text{ has degenerate PDF.)}} \mathbb{1}_{\{\|\tau - \tilde{\nu}\|_1 \leq t\}} d\tau$$

$$\leq \left(\sup_{\tau \in \mathbb{B}^{n-1}} P_{\tilde{\pi}}(\tau) \right) \cdot \underbrace{\text{vol}(\{x \in \mathbb{B}^{n-1} : \|x\|_1 \leq t\})}_{\text{volume of } \ell^1\text{-ball in } \mathbb{B}^{n-1} \text{ with radius } t = \frac{2^{n-1}}{(n-1)!} t^{n-1}}$$

$$= \frac{2^{n-1}}{(n-1)!} t^{n-1} \cdot \sup_{\tau \in \mathbb{B}^{n-1}} P_{\tilde{\pi}}(\tau)$$

Lemma:
 $\sup_{\tau \in \mathbb{B}^{n-1}} P_{\tilde{\pi}}(\tau) \leq \frac{n^{n-1}}{(1-\delta)^n}$

(Follows from explicit calculation of PDF $P_{\tilde{\pi}}$ from P_{α_n} based on change-of-var. formula.)

$$\leq \frac{2^{n-1}}{(n-1)!} t^{n-1} \cdot \frac{n^{n-1}}{(1-\delta)^n}$$

$$= \frac{1}{2} \cdot \frac{n^n}{n!} \left(\frac{2}{1-\delta}\right)^n t^{n-1}$$

$$\leq \frac{1}{5\sqrt{n}} \left(\frac{2e}{1-\delta}\right)^n t^{n-1} \left\} \text{Stirling's formula: } n! \geq \frac{5}{2} \sqrt{n} \frac{n^n}{e^n}.$$

$$\leq \left(\frac{2e}{1-\delta}\right)^n t^{n-1}.$$

This completes the proof. ▀

• Theorem: For any $\varepsilon > 0$,

(Lower Bound)

$$\inf_{\hat{\pi}} \sup_{P_{\alpha} \in \mathcal{P}_{\delta}(\sigma)} \mathbb{E}[\|\hat{\pi} - \pi\|_1] \geq \left(\frac{\varepsilon}{4+2\varepsilon}\right) \frac{1}{n^{\frac{1}{2}+\varepsilon}}$$

for all $n \geq 2$ sufficiently large (depending on ε, δ).

↑ $n \uparrow$ as $\varepsilon \downarrow 0$, $n \uparrow$ as $\delta \downarrow 0$

- Remark: For any $\varepsilon > 0$, and all $n \in \mathbb{N}$ sufficiently large:

$$\Omega\left(\frac{1}{n^{\frac{1}{2}+\varepsilon}}\right) \leq \inf_{\hat{\pi}} \sup_{P_{\alpha} \in \mathcal{P}_{\delta}(\sigma)} \mathbb{E}[\|\hat{\pi} - \pi\|_1] \leq O\left(\frac{1}{\sqrt{n}}\right).$$

[continued.]

- Proof: Fix any $\varepsilon > 0$. Using previous lemmata:

$$\begin{aligned}
 \inf_{\hat{\pi}} \sup_{P_{\alpha} \in \mathcal{B}_X(\sigma)} \mathbb{E}[\|\hat{\pi} - \pi\|_1] &\geq \inf_{\hat{\pi}} \mathbb{E}[\|\hat{\pi} - \pi\|_1] \\
 &\quad \uparrow P_{\alpha} = \text{Uniform}([s, 1]) \\
 &\geq \sup_{t > 0} t \left(1 - \frac{I(\alpha^n; Z) + \log(2)}{\log(1/\mathcal{L}(t))} \right) \quad [\text{Generalized Fano Lemma}] \\
 &\geq \sup_{t > 0} t \left(1 - \frac{\frac{1}{2}n\log(n) + \frac{(1-\delta)^2}{8\delta^2}(2+\delta+\frac{1}{\delta})n + \log(2)}{(n-1)\log(1/t) - \log(2e/(1-\delta))n} \right) \quad [\text{Upper Bounds on } I(\alpha^n; Z) \text{ and } \mathcal{L}(t)] \\
 &\geq \sup_{t > 0} t \left(1 - \frac{1 + O\left(\frac{1}{\log(n)}\right)}{\frac{2(n-1)\log(1/t)}{n\log(n)} - O\left(\frac{1}{\log(n)}\right)} \right) \\
 &\quad \downarrow \text{Set } t = \frac{1}{n^{\frac{1}{2}+\varepsilon}} \\
 &\geq \frac{1}{n^{\frac{1}{2}+\varepsilon}} \left(1 - \frac{1 + O(1/\log(n))}{(1+2\varepsilon)(1-\frac{1}{n}) - O(1/\log(n))} \right) \\
 &\quad \rightarrow 1 - \frac{1}{1+2\varepsilon} = \frac{2\varepsilon}{1+2\varepsilon} \text{ as } n \rightarrow \infty \\
 &\geq \frac{1}{n^{\frac{1}{2}+\varepsilon}} \left(\frac{\varepsilon}{4+2\varepsilon} \right) \text{ for } n \geq 2 \text{ sufficiently large.}
 \end{aligned}$$

This completes the proof. ▀

~ ● ~